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Characterization Theorems for Modular and Distributive Soft Lattices

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Abstract: Soft set theory was introduced by Molodtsov in 1999 as a mathematical tool for dealing with problems that contain uncertainty. Faruk Karaaslan et al.[6] defined the concept of soft lattices, modular soft lattices and distributive soft lattice over a collection of soft sets. In this paper, we define the concept of principle of duality in soft lattices and discuss some related properties of modular and distributive soft lattices. We also establish characterization theorems for modular and distributives soft sublattices.

Key words: Soft sets, soft lattices, soft sublattices, dual soft lattices, modular soft lattices, distributive soft lattices.

I. INTRODUCTION

Soft set theory was introduced by Molodtsov [9] in 1999 as a mathematical tool for dealing with uncertainty. Maji et al.[8] defined some operations on soft sets and proved related properties. Irfan Ali et al.[5] studied some new operations on soft sets. Li [7], Nagarajan et al.[10] defined the soft lattices using soft sets. Faruk Karaaslan et al.[6] defined the concept of soft lattices over a collection of soft sets by using the operations of soft sets defined by Cagman et al.[1]. In this paper, we define the concept of principle of duality in soft lattices and discuss some related properties of modular and distributive soft lattices. We also illustrate them with some examples. In addition, we establish characterization theorems for modular and distributive soft lattices by their soft sublattices.

II. SOME CONCEPTS IN SOFT SETS AND SOFT LATTICES

In this section, we have presented the basic definitions and results of soft set theory [9, 6] and soft lattices [6] for subsequent discussions. Throughout this work, U refers to the initial universe, P(U) is the power set of U, E is a set of parameters and $A \subseteq E$.

Definition 2.1 A function $f_A: E \to P(U)$ such that $f_A(x) = \emptyset$ if $x \notin A$, is called a soft set over U. The set of all soft sets over U is denoted by S(U).

Definition 2.2 Let $f_A \in S(U)$. If $f_A(x) = \emptyset$ for all $x \in E$, then f_A is called an empty soft set, denoted by f_{\emptyset} . If $f_A(x) = U$ for all $x \in A$, then f_A is called A-universal soft set denoted by $f_{\tilde{A}}$. If A = E then the

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A -universal soft set is called universal soft set denoted by $f_{\tilde{F}}$.

Definition 2.3 Let $f_A, f_B \in S(U)$. Then, f_A is a soft subset of f_B , denoted by $f_A \cong f_B$, if $f_A(x) \subseteq f_B(x)$ for all $x \in E$. f_A and f_B are equal, denoted by $f_A = f_B$, if and only if $f_A(x) = f_B(x)$ for all $x \in E$.

Definition 2.4 Let $f_A \in S(U)$. Then, soft complement of f_A is defined by $f_A^{\tilde{c}} = f_{A^{\tilde{c}}}$ such that $f_{A^{\tilde{c}}}(x) = f_A^{\tilde{c}}(x) = U \setminus f_A(x)$ for all $x \in E$.

Definition 2.5 Let $f_A, f_B \in S(U)$. Then, soft union of f_A and f_B is defined by $f_A \widetilde{\cup} f_B = f_{A\widetilde{\cup}B}$ such that $f_{A\widetilde{\cup}B}(x) = f_A(x) \cup f_B(x)$ for all $x \in E$.

Definition 2.6 Let $f_A, f_B \in S(U)$. Then, soft intersection of f_A and f_B is defined by $f_A \cap f_B = f_{A \cap B}$ such that $f_{A \cap B}(x) = f_A(x) \cap f_B(x)$ for all $x \in E$.

Definition 2.7 Let $L \subseteq S(U)$, and \lor and \land be two binary operations on L. If the set L together with \lor and \land satisfies the following conditions then (L,\lor,\land) is called a soft lattice.

 $\begin{array}{ll} 1. \quad f_A \lor f_B = f_B \lor f_A \ \text{and} \ f_A \land f_B = f_B \land f_A \ \text{for all} \ f_A, f_B \in L. \\ 2. \quad f_A \lor (f_B \lor f_C) = (f_A \lor f_B) \lor f_C \ \text{and} \ f_A \land (f_B \land f_C) = (f_A \land f_B) \land f_C \ \text{for all} \\ f_A, f_B, f_C \in L \ . \\ 3. \quad f_A \lor (f_A \land f_B) = f_A \ \text{and} \ f_A \land (f_A \lor f_B) = f_A \ \text{for all} \ f_A, f_B \in L \ . \end{array}$

Theorem 2.8 Let (L, \lor, \land) be a soft lattice and $f_A, f_B \in L$. Then, $f_A \land f_B = f_A \Leftrightarrow f_A \lor f_B = f_B$.

Theorem 2.9 Let (L, \lor, \land) be a soft lattice and $f_A, f_B \in L$. Then a relation \leq that is defined by $f_A \leq f_B \Leftrightarrow f_A \land f_B = f_A$ or $f_A \lor f_B = f_B$ is an ordering relation on L.

Lemma 2.10 Let (L, \lor, \land) be a soft lattice and $f_A, f_B \in L$. Then, $f_A \lor f_B$ and $f_A \land f_B$ are the least upper bound and the greatest lower bound of f_A and f_B , respectively.

Theorem 2.11 A soft lattice (L, \lor, \land) is a poset.

Theorem 2.12 Let $L \subseteq S(U)$. Then, the algebraic structure (L, \lor, \land, \leq) is a soft lattice.

Note 2.13 According to the theorem 2.12, a soft lattice (L,\vee,\wedge) has the same character with (L,\vee,\wedge,\leq) . Therefore, we shall identify any soft lattice (L,\vee,\wedge) with (L,\vee,\wedge,\leq) and we use these two concepts

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interchangeably.

Definition 2.14 Let (L, \lor, \land, \leq) be a soft lattice and $S \subseteq L$. If S is a soft lattice with the operations of L, then S is called a soft sublattice of L.

Theorem 2.15 Let (L, \lor, \land, \leq) be a soft lattice and $S \subseteq L$. If $f_A \lor f_B \in S$ and $f_A \land f_B \in S$ for all $f_A, f_B \in S$, then S is a soft sublattice.

III. MODULAR AND DISTRIBUTIVE SOFT LATTICES

In this section, we introduce the concept of duality principle and give the definition of modular soft lattice and distributive soft lattice. Also we discuss several related properties and characterization theorems on modular and distributive soft lattices. Throughout this section, U refers to the initial universe. The set of all soft sets over U is denoted by S(U). $L \subseteq S(U)$ and (L, \lor, \land, \leq) is a soft lattice.

Definition 3.1 Duality Principle

Let (L, \lor, \land, \le) be a soft lattice. If we define a relation \le in L for all $f_A, f_B \in L$, $f_A \le f_B \Leftrightarrow f_B \le f_A$ in (L, \lor, \land, \le) , then \le is also a partial ordering on L. The partial ordering \le is called the reversal partial ordering of \le .

Clearly, for all $f_A, f_B \in L$, l.u.b $\{f_A, f_B\}$ in (L, \lor, \land, \leq) =g.l.b $\{f_A, f_B\}$ in (L, \lor, \land, \leq) and g.l.b $\{f_A, f_B\}$ in (L, \lor, \land, \leq) = l.u.b $\{f_A, f_B\}$ in (L, \lor, \land, \leq) . The operations \lor and \land with respect to \leq in L are the \land and \lor with respect to \leq' in L. As $f_A \leq' f_B$ in (L, \lor, \land, \leq') means $f_B \leq f_A$ in (L, \lor, \land, \leq) , any formula involving \lor, \land, \leq in (L, \lor, \land, \leq) remains valid in (L, \lor, \land, \leq') if \leq, \lor, \land are replaced by \geq, \land, \lor

Remark 3.2 In any soft lattice (L, \lor, \land, \le) for all $f_A, f_B \in L$, $f_A \le f_B \Longrightarrow f_A \lor f_B = f_B$ is valid. Hence its dual formula $f_A \ge f_B \Longrightarrow f_A \land f_B = f_B$ is valid.

Definition 3.3 Let (L, \lor, \land, \le) be a soft lattice. Then L is said to be a modular soft lattice if $f_A \le f_C \Longrightarrow f_A \lor (f_B \land f_C) = (f_A \lor f_B) \land f_C$ for all $f_A, f_B, f_C \in L$. or equivalently,

$$f_A \ge f_C \Longrightarrow f_A \land (f_B \lor f_C) = (f_A \land f_B) \lor f_C \text{ for all } f_A, f_B, f_C \in L$$

Remark 3.4 To show that the soft lattice L is a Modular soft lattice it is enough to prove that $f_A \leq f_C \Longrightarrow (f_A \lor f_B) \land f_C \leq f_A \lor (f_B \land f_C)$ for all $f_A, f_B, f_C \in L$, as the other inequality $f_A \lor (f_B \land f_C) \leq (f_A \lor f_B) \land f_C$ for all $f_A, f_B, f_C \in L$ is always true for all soft lattices.

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 $D = \{e_1, e_2, e_3, e_4\}$ where $A, B, C, D \subseteq E$ and $L = \{f_{\emptyset}, f_A, f_B, f_C, f_D\}$. Then $L \subseteq S(U)$ is a soft lattice with the operations $\tilde{\cup}$ and $\tilde{\cap}$. Assume that,

$$\begin{split} f_{\varnothing} &= \varnothing \\ f_A &= \{(e_1, \{u_1, u_2\})\} \\ f_B &= \{(e_2, \{u_3, u_4\})\} \\ f_C &= \{(e_3, \{u_5, u_6\}), (e_4, \{u_1\})\} \\ f_D &= \{(e_1, \{u_1, u_2\}), (e_2, \{u_3, u_4\}), (e_3, \{u_5, u_6\}), (e_4, \{u_1, u_2\})\} \end{split}$$

 $(L, \widetilde{\bigcirc}, \widetilde{\bigcirc}, \widetilde{\subseteq})$ is a Modular soft lattice. The Hasse diagram of it appears in figure 1. This Modular soft lattice is denoted by SM_3 .



Figure 1: Modular Soft Lattice (SM₃)

Example 3.6 Let $U = \{u_1, u_2, u_3, u_4, u_5\}, E = \{e_1, e_2, e_3\}, A = \{e_1\}, B = \{e_2\}, C = \{e_1, e_3\}, D = \{e_1, e_2, e_3\}$ where $A, B, C, D \subseteq E$ and $L = \{f_{\emptyset}, f_A, f_B, f_C, f_D\}$. Then $L \subseteq S(U)$ is a soft lattice with the operations $\tilde{\cup}$ and $\tilde{\cap}$. Assume that,

$$\begin{split} f_{\varnothing} &= \varnothing \\ f_{A} &= \{(e_{1}, \{u_{1}\})\} \\ f_{B} &= \{(e_{2}, \{u_{2}, u_{3}\})\} \\ f_{C} &= \{(e_{1}, \{u_{1}, u_{2}\}), (e_{3}, \{u_{3}, u_{4}, u_{5}\})\} \\ f_{D} &= \{(e_{1}, \{u_{1}, u_{2}\}), (e_{2}, \{u_{2}, u_{3}\}), (e_{3}, \{u_{3}, u_{4}, u_{5}\}\} \end{split}$$

 $(L, \widetilde{\bigcirc}, \widetilde{\bigcirc}, \widetilde{\subseteq})$ is a Non-modular soft lattice. The Hasse diagram of it appears in figure 2. This Non-modular soft lattice is denoted by SN_5 .



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Figure 2: Non-Modular Soft Lattice (SN₅)

Definition 3.7 Let (L, \lor, \land, \leq) be a soft lattice. Then L is said to be a distributive soft lattice if for all $f_A, f_B, f_C \in L$

1.
$$f_A \lor (f_B \land f_C) = (f_A \lor f_B) \land (f_A \lor f_C)$$

2. $f_A \land (f_B \lor f_C) = (f_A \land f_B) \lor (f_A \land f_C)$

[These equations are called distributive equations.]

Remark 3.8 Because of duality principle, a soft lattice L is said to be a distributive soft lattice if any one of the distributive equations is true for all $f_A, f_B, f_C \in L$.

Remark 3.9 To show that the soft lattice L is a distributive soft lattice it is enough to prove that, $f_A \wedge (f_B \vee f_C) \leq (f_A \wedge f_B) \vee (f_A \wedge f_C)$ for all $f_A, f_B, f_C \in L$ as the other inequality $(f_A \wedge f_B) \vee (f_A \wedge f_C) \leq f_A \wedge (f_B \vee f_C)$ is always true for all $f_A, f_B, f_C \in L$.

Example 3.10 Let $U = \{u_1, u_2, u_3, u_4, u_5, u_6\}, E = \{e_1, e_2, e_3\}, A = \{e_1\}, B = \{e_1, e_2\}, C = \{e_1, e_3\}, D = \{e_1, e_2, e_3\}$ where $A, B, C, D \subseteq E$ and $L = \{f_{\emptyset}, f_A, f_B, f_C, f_D\}$. Then $L \subseteq S(U)$ is a soft lattice with the operations $\tilde{\cup}$ and $\tilde{\cap}$. Assume that

$$\begin{split} f_{\varnothing} &= \varnothing \\ f_A &= \{(e_1, \{u_1, u_2\})\} \\ f_B &= \{(e_1, \{u_1, u_2, u_3\}), (e_2, \{u_4, u_6\})\} \end{split}$$

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$$\begin{aligned} f_C &= \{(e_1, \{u_1, u_2\}), (e_3, \{u_5, u_6\})\} \\ f_D &= \{(e_1, \{u_1, u_2, u_3\}), (e_2, \{u_4, u_6\}), (e_3, \{u_5, u_6\})\} \\ (L, \widetilde{\cup}, \widetilde{\cap}, \underline{\widetilde{\frown}}) & \text{ is a distributive soft lattice. The Hasse diagram of it appears in figure 3.} \end{aligned}$$



Figure 3: Distributive Soft Lattice

Lemma 3.11 The soft sublattice (dual) of a modular soft lattice is modular.

Proof. Let L be a modular soft lattice. Then $f_A \leq f_C \Rightarrow f_A \vee (f_B \wedge f_C) = (f_A \vee f_B) \wedge f_C$ for all $f_A, f_B, f_C \in L$. Let S be a soft sublattice of L.Take $f_A, f_B, f_C \in S$. Since $S \subseteq L$, $f_A, f_B, f_C \in L$. Therefore, $f_A \leq f_C \Rightarrow f_A \vee (f_B \wedge f_C) = (f_A \vee f_B) \wedge f_C$ is true for all $f_A, f_B, f_C \in S$. Hence the soft sublattice S is modular soft lattice. Therefore, the soft sublattice of a modular soft lattice is modular.

Let L be a modular soft lattice. Then, $f_A \leq f_C \Rightarrow f_A \lor (f_B \land f_C) = (f_A \lor f_B) \land f_C$ for all $f_A, f_B, f_C \in L$. The dual of L is $f_A \geq f_C \Rightarrow f_A \land (f_B \lor f_C) = (f_A \land f_B) \lor f_C$ for all $f_A, f_B, f_C \in L$. Therefore, the dual of the modular soft lattice is modular.

Lemma 3.12 A soft lattice L is modular if and only if every triplet of elements $f_A, f_B, f_C(f_A \le f_C)$ in L has a median.

(or) A Soft lattice *L* is modular iff for all $f_A, f_B, f_C \in L$ with $f_A \leq f_C$, $(f_A \wedge f_B) \vee (f_B \wedge f_C) \vee (f_C \wedge f_A) = (f_A \vee f_B) \wedge (f_B \vee f_C) \wedge (f_C \vee f_A).$

Proof. Suppose that L is a soft lattice in which every triplet of elements $f_A, f_B, f_C(f_A \le f_C)$ has a median.



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To prove that L is a modular soft lattice.

From (1) and (2),we get, $f_A \lor (f_B \land f_C) = (f_A \lor f_B) \land f_C$

Hence L is a modular soft lattice.

Conversely, suppose that L is a modular soft lattice.

Then, $f_A \leq f_C \Longrightarrow f_A \lor (f_B \land f_C) = (f_A \lor f_B) \land f_C$

To prove that every triplet of elements f_A, f_B, f_C $(f_A \leq f_C)$ in L has a median.

From (3) and (4), we get, $(f_A \wedge f_B) \vee (f_B \wedge f_C) \vee (f_C \wedge f_A) = (f_A \vee f_B) \wedge (f_B \vee f_C) \wedge (f_C \vee f_A)$

Theorem 3.13 A soft lattice L is modular if and only if it has no soft sublattice isomorphic to SN_5 .

Proof. Let L be a modular soft lattice. Then every soft sublattice of a modular soft lattice is modular.(1)

Consider the soft lattice SN_5 in figure 4.





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Figure 4.

Here $f_A \leq f_C$, $f_A \vee (f_B \wedge f_C) = f_A \vee f_{\varnothing} = f_A$ while $(f_A \vee f_B) \wedge f_C = f_D \wedge f_C = f_C$. So $f_A \vee (f_B \wedge f_C) \neq (f_A \vee f_B) \wedge f_C$. Hence, SN_5 is not a modular soft lattice. Since every sublattice of a modular soft lattice is modular, L cannot have a soft sublattice isomorphic to SN_5 . Conversely, assume that L has no soft sublattice isomorphic to SN_5 . To prove, L is a modular soft lattice. That is, $f_A \leq f_C \Rightarrow f_A \vee (f_B \wedge f_C) = (f_A \vee f_B) \wedge f_C$ for all $f_A, f_B, f_C \in L$. Suppose L is not a modular soft lattice, then there exist elements $f_A, f_B, f_C \in L$ such that $f_A \leq f_C$ and $f_A \vee (f_B \wedge f_C) < (f_A \vee f_B) \wedge f_C$.

Let
$$f_P = f_B \wedge f_C$$
, $f_X = f_A \vee (f_B \wedge f_C)$, $f_Y = f_B$, $f_Z = (f_A \vee f_B) \wedge f_C$, $f_Q = (f_A \vee f_B)$
Let $S = \{f_P, f_X, f_Y, f_Z, f_Q\}$

We claim that S is a soft sublattice of L isomorphic to SN_5 .

By assumption, we have $f_P \leq f_X \leq f_Z \leq f_Q$ and $f_P \leq f_Y \leq f_Q$.

 $\begin{array}{l} 1. \ \ f_P \leq f_X \wedge f_Y \leq f_Z \wedge f_Y = (f_A \vee f_B) \wedge f_C \wedge f_B = ((f_A \vee f_B) \wedge f_B) \wedge f_C = f_B \wedge f_C = f_P. \\ \text{So} \ \ f_X \wedge f_Y = f_Z \wedge f_Y = f_P. \\ 2. \ \ f_Q \geq f_Z \vee f_Y \geq f_X \vee f_Y = f_A \vee (f_B \wedge f_C) \vee f_B = f_A \vee ((f_B \wedge f_C) \vee f_B) = f_A \vee f_B = f_Q. \\ \text{So} \ \ f_Z \vee f_Y = f_X \vee f_Y = f_O. \end{array}$

Therefore, f_P and f_O are the bound elements of the soft sublattice S.

It remains to prove that these five elements f_P, f_X, f_Y, f_Z, f_O are all distinct. (ie),to prove





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 $f_P \neq f_Y, f_Q \neq f_Y, f_P \neq f_X, f_Q \neq f_Z, f_X \neq f_Y, f_Z \neq f_Y$ Suppose $f_P = f_Y \Longrightarrow f_X \land f_Y = f_Y \Longrightarrow f_X \lor f_Y = f_X \implies f_Q = f_X$ which is impossible for $f_P \leq f_X \leq f_Z \leq f_Q$. Therefore, $f_P \neq f_Y$ Suppose $f_Q = f_Z \Longrightarrow f_Z \lor f_Y = f_Z \implies f_Z \land f_Y = f_Y \implies f_P = f_Y$ which is a contradiction. Therefore, $f_Q \neq f_Z$ Suppose $f_X = f_Y$. Here $f_P = f_X \land f_Y \implies f_P = f_Y \land f_Y \implies f_P = f_Y$ which is a contradiction. Therefore, $f_X \neq f_Y$ Suppose $f_X = f_Y$. Here $f_P = f_X \land f_Y \implies f_P = f_Y \land f_Y \implies f_P = f_Y$ which is a contradiction. Therefore, $f_X \neq f_Y$ By dualizing the above statement, we get $f_Q \neq f_Y, f_P \neq f_X, f_Z \neq f_Y$.

These observations lead to the representation of $S = \{f_P, f_X, f_Y, f_Z, f_O\}$ by a Hasse diagram given in figure 5.





Therefore, S is a soft sublattice of L isomorphic to SN_5 , which is a contradiction. Hence L is a modular soft lattice.

Lemma 3.14 *The soft sublattice (dual) of a distributive soft lattice is distributive.*

Proof. Let L be a distributive soft lattice. Then $f_A \wedge (f_B \vee f_C) = (f_A \wedge f_B) \vee (f_A \wedge f_C)$ for all $f_A, f_B, f_C \in L$. Let S be a soft sublattice of L. Take $f_A, f_B, f_C \in S$. Since $S \subseteq L$, $f_A, f_B, f_C \in L$. Therefore, $f_A \wedge (f_B \vee f_C) = (f_A \wedge f_B) \vee (f_A \wedge f_C)$ is true for all $f_A, f_B, f_C \in S$. Hence the soft sublattice S is a distributive soft lattice. Therefore, the soft sublattice of a distributive soft lattice is distributive.

Let L be a distributive soft lattice. Then $f_A \wedge (f_B \vee f_C) = (f_A \wedge f_B) \vee (f_A \wedge f_C)$ for all $f_A, f_B, f_C \in L$. The dual of L is $f_A \vee (f_B \wedge f_C) = (f_A \vee f_B) \wedge (f_A \vee f_C)$ for all $f_A, f_B, f_C \in L$. Therefore, the dual of

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the distributive soft lattice is distributive.

Lemma 3.15 A soft lattice L is distributive if and only if every one of its triplets has a median. (or) A soft lattice L is distributive $\Leftrightarrow (f_A \lor f_B) \land (f_B \lor f_C) \land (f_C \lor f_A) = (f_A \land f_B) \lor (f_B \land f_C) \lor (f_C \land f_A)$ for all $f_A, f_B, f_C \in L$.

$$\begin{aligned} \text{Proof. Let } L \text{ be a distributive soft lattice.} \\ \text{Let } f_A, f_B, f_C \in L \\ (f_A \lor f_B) \land (f_B \lor f_C) \land (f_C \lor f_A) \\ &= \{(f_A \lor f_B) \land (f_B \lor f_C)\} \land (f_C \lor f_A) \\ &= \{((f_A \lor f_B) \land (f_B \lor f_C)) \land (f_C)\} \lor \{((f_A \lor f_B) \land (f_B \lor f_C)) \land (f_A)\} \\ &= \{(f_A \lor f_B) \land ((f_B \lor f_C) \land (f_C))\} \lor \{(f_A \land (f_A \lor f_B)) \land (f_B \lor f_C)\} \\ &= \{(f_A \lor f_B) \land f_C\} \lor \{f_A \land (f_B \lor f_C)\} \\ &= \{(f_A \land f_C) \lor (f_B \land f_C)\} \lor \{(f_A \land f_B) \lor (f_A \land f_C)\} \\ &= \{(f_A \land f_B) \lor (f_B \land f_C) \lor (f_C \land f_A) \lor (f_C \land f_A) \\ &= (f_A \land f_B) \lor (f_B \land f_C) \lor (f_C \land f_A). \end{aligned}$$

Therefore, if L is a distributive soft lattice then every triplet of elements in L has a median. Conversely, Suppose that L is a soft lattice in which every triplet of elements has a median. To prove that L is a distributive soft lattice.

Let
$$f_A, f_B, f_C \in L$$
.

$$\begin{aligned} f_A \wedge (f_B \vee f_C) &= (f_A \wedge (f_A \vee f_C)) \wedge (f_B \vee f_C) \\ &= f_A \wedge (f_A \vee f_C) \wedge (f_B \vee f_C) \\ &= (f_A \wedge (f_A \vee f_B)) \wedge (f_A \vee f_C) \wedge (f_B \vee f_C)) \\ &= f_A \wedge ((f_A \wedge f_B) \vee (f_A \wedge f_C) \vee (f_B \wedge f_C)) \\ &= f_A \wedge ((f_B \wedge f_C) \vee (f_A \wedge f_B) \vee (f_C \wedge f_A)) \end{aligned}$$
Since, L is modular, $f_A \geq (f_A \wedge f_B) \vee (f_A \wedge f_C)$, we get
$$\begin{aligned} f_A \wedge (f_B \vee f_C) &= \{f_A \wedge (f_B \wedge f_C)\} \vee \{(f_A \wedge f_B) \vee (f_C \wedge f_A)\} \\ &= \{(f_A \wedge f_B \wedge f_C) \vee (f_A \wedge f_B)\} \vee (f_A \wedge f_C) \\ &= (f_A \wedge f_B) \vee (f_A \wedge f_C) \end{aligned}$$

Hence L is a distributive soft lattice.

Theorem 3.16 A soft lattice L is distributive if and only if it has no soft sublattice isomorphic to SN_5 or SM_3 . *Proof.* Let L be a distributive soft lattice. The soft lattice SN_5 is not a modular soft lattice and so it cannot be a

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distributive soft lattice. Consider SM_3 in figure 6.



Figure 6.

 $f_A \vee (f_B \wedge f_C) = f_A \vee f_{\varnothing} = f_A$; $(f_A \vee f_B) \wedge (f_A \vee f_C) = f_D \vee f_D = f_D$ Therefore, $f_A \vee (f_B \wedge f_C) \neq (f_A \vee f_B) \wedge (f_A \vee f_C)$. Hence SM_3 is not a distributive soft lattice. If L is distributive soft lattice then every soft sublattice of L is also distributive. So L cannot have a soft

sublattice isomorphic to SN_5 or SM_3 . Conversely, Suppose that L is a soft lattice which has no soft sublattice isomorphic to SN_5 or SM_3 . By theorem 3.13, it is enough if we prove that, a modular soft lattice is non- distributive then it has a soft sublattice isomorphic to SM_3 . Let L be a modular non-distributive soft lattice. Then L has a triplet f_P, f_Q, f_R of elements such that $(f_P \wedge f_Q) \vee (f_Q \wedge f_R) \vee (f_R \wedge f_P) < (f_P \vee f_Q) \wedge (f_Q \vee f_R) \wedge (f_R \vee f_P)$ Let $f_X = (f_P \wedge f_Q) \vee (f_Q \wedge f_R) \vee (f_R \wedge f_P)$ $f_Y = (f_P \vee f_Q) \wedge (f_Q \vee f_R) \wedge (f_R \vee f_P)$.

Then $f_X \leq f_Y$ and L is modular soft lattice. Let $f_A = f_X \vee (f_P \wedge f_Y) = (f_X \vee f_P) \wedge f_Y$, $f_B = f_X \vee (f_Q \wedge f_Y) = (f_X \vee f_Q) \wedge f_Y$ $f_C = f_X \vee (f_R \wedge f_Y) = (f_X \vee f_R) \wedge f_Y$

Let $S = \{f_X, f_Y, f_A, f_B, f_C\}$. We claim that S is a soft sublattice of L isomorphic to SM_3 . To prove $f_A \wedge f_B = f_B \wedge f_C = f_C \wedge f_A = f_X$ and $f_A \vee f_B = f_B \vee f_C = f_C \vee f_A = f_Y$ Here, $f_P \wedge f_Y = f_P \wedge \{(f_P \vee f_Q) \wedge (f_Q \vee f_R) \wedge (f_R \vee f_P)\} = f_P \wedge (f_Q \vee f_R)$

$$\begin{split} f_{\mathcal{Q}} \wedge f_{Y} &= f_{\mathcal{Q}} \wedge \{(f_{P} \vee f_{Q}) \wedge (f_{\mathcal{Q}} \vee f_{R}) \wedge (f_{R} \vee f_{P})\} = f_{\mathcal{Q}} \wedge (f_{R} \vee f_{P}) \\ \text{Now,} \ f_{A} \vee f_{B} &= f_{X} \vee \{(f_{P} \wedge (f_{\mathcal{Q}} \vee f_{R})) \vee (f_{\mathcal{Q}} \wedge (f_{R} \vee f_{P}))\} \dots \dots \dots \dots \dots (1) \\ \text{since} \ f_{P} \wedge (f_{\mathcal{Q}} \vee f_{R}) \leq f_{P} \leq f_{R} \vee f_{P} \\ (f_{P} \wedge (f_{\mathcal{Q}} \vee f_{R})) \vee (f_{\mathcal{Q}} \wedge (f_{R} \vee f_{P})) = ((f_{P} \wedge (f_{\mathcal{Q}} \vee f_{R})) \vee f_{\mathcal{Q}}) \wedge (f_{R} \vee f_{P}) \\ &= (f_{\mathcal{Q}} \vee ((f_{P} \wedge (f_{\mathcal{Q}} \vee f_{R})) \wedge (f_{R} \vee f_{P})) \\ \end{split}$$

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$$= (f_Q \lor f_P) \land (f_Q \lor f_R) \land (f_R \lor f_P)$$

Therefore,(1) becomes $f_A \lor f_B = f_X \lor ((f_Q \lor f_P) \land (f_Q \lor f_R) \land (f_R \lor f_P))$
 $= f_X \lor f_Y = f_Y$
Hence $f_A \lor f_B = f_Y$. Similarly we prove, $f_B \lor f_C = f_C \lor f_A = f_Y$
By the principle of duality, we have $f_A \land f_B = f_X$ and $f_B \land f_C = f_C \land f_A = f_X$
Therefore, f_X and f_Y are the bound elements of the soft sublattice S .
It remains to prove that these five elements f_X, f_A, f_B, f_C, f_Y are all distinct. That is to prove
 $f_X \neq f_A, f_X \neq f_B, f_X \neq f_C, f_Y \neq f_A, f_Y \neq f_B, f_Y \neq f_C$ and $f_A \neq f_B, f_B \neq f_C, f_C \neq f_A$
Suppose $f_X = f_A$
 $\Rightarrow f_A \land f_B = f_A$
 $\Rightarrow f_A \land f_B = f_B$
 $\Rightarrow f_Y = f_B$
Now, $f_A \land f_C = f_A$
 $\Rightarrow f_A \lor f_C = f_C$
Therefore, $f_Y = f_Y \land f_Y = f_B \land f_C = f_X$ which is a contradiction. Hence $f_X \neq f_A$.
Similarly, we have $f_X \neq f_B, f_X \neq f_C$.
If $f_A = f_B$ then $f_X = f_A \land f_A = f_A$ which is a contradiction. Hence $f_A \neq f_B$.
Similarly we have $f_B \neq f_A$, $f_Y \neq f_B, f_Y \neq f_C$.
If $f_A = f_B$ then $f_X = f_A \land f_A = f_A$ which is a contradiction. Hence $f_A \neq f_B$.
Similarly we have $f_B \neq f_C, f_C \neq f_A$.

These observations lead to the representation of $S = \{f_X, f_A, f_B, f_C, f_Y\}$ by a Hasse diagram given in figure 7.





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Therefore, S is a soft sublattice of L isomorphic to SM_3 . Thus a modular non-distributive soft lattice has an isomorphic copy of SM_3 . Therefore, if L has no soft sublattice isomorphic to SN_5 or SM_3 then L is a distributive soft lattice.

IV.CONCLUSION

In this paper, we have given characterization theorems for modular soft lattices and distributive soft lattices. We are studying about these soft lattices and are expected to give some more results.

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