

(An ISO 3297: 2007 Certified Organization) Vol. 4, Issue 12, December 2015

Implementation of Binary Extension Field for Quad Core Cryptoprocessor on FPGA

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ABSTRACT: This paper is devoted to the design of Quad core crypto processor for executing binary extension field instructions. The proposed design is specifically optimized for Field programmable gate array (FPGA) platform. Combination of two different field (prime field GF(P) and Binary extension field GF(2m)) instructions execution is analyzed. Quad core will execute four instruction at a same time. The design is implemented in Spartan 3E, virtex4 and virtex5. The performance results between them are compared. The implementation result shows the execution of parallelism using dual field instructions.

KEYWORDS: Binary extension field, Cryptoprocessor, FPGA, Prime field

I. INTRODUCTION

The term cryptography represents the encryption of data. For the secured data communication cryptography is essential one. Cryptography is used now a day in a variety of different applications. A secure crypto processor is a dedicated microprocessor chip for carrying out cryptographic operation. Every application has its own design criteria and raises special requirements for hardware designs. The most fundamental decision concerning future hardware designs is whether to use a binary-extension field or a prime field as basis of the used crypto processor. Only a few papers compared binary and prime fields in hardware. Dual-Field Arithmetic Unit for GF(p) and GF(2m) is designed by Johannes Wolkerstorfer [1]. Jun-Hong Chen et al. designed the high performance unified field reconfigurable cryptoprocessor. The cryptography operations involving the integer value is carry over by GF(p). The binary extension fields GF(2m), where elements can be represented as polynomials instead of integers. Binary fields GF(2m) are considered advantageous for hardware solutions because addition and modular reduction of polynomials are somewhat easier than those of integers. We will present a dualfield arithmetic unit that is capable to calculate the operations in both types of fields: GF(p) and GF(2m) A reconfigurable multicore cryptoprocessor for multi channel communication systems [3] result shows the increase in speed of execution. For increasing the speed of execution this paper is designing the quad core crypto processor. The quad core will execute the instruction concurrently. The parallel execution will increase the speed of execution. The parallelization of the hardware design is a significant ongoing topic of research. At the same time for decreasing the power consumption this design is implemented in the FPGA. Santosh Ghosh et al. design the secure dual core cryptoprocessor using the prime field



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instructions [4]. This paper is concentrate to design a quad core crypto processor used to execute both prime and binary extension field instructions.

II. LITERATURE SURVEY

Hassan A. Youness, Mahmoud Khaled, and Mohamed Moness, "Quad-core MPSoC architecture For PID-based Embedded Control Systems" The main contribution of this paper is a throughput-oriented high-performance low-cost digital PID controller.

Michael Grand1, Lilian Bossuet2, Guy Gogniat3, Bertrand Le Gal1,"A Reconfigurable Multi-core Cryptoprocessor For Multi-channel Communication Systems" The proposed MCCP can reach a maximum throughput of 1.7 Gbps at 190 MHz with several AES cipher modes. It uses about 4000 slices on a Virtex 4 FPGA.

Jun-Hong Chen, Ming-Der Shieh, Member, IEEE, *and* Wen-Ching Lin, "A High Performance Unified-field Reconfigurable Cryptographic processor" Rivest –Shamir- Adleman(RSA), Elliptic Curve Crytography(ECC) - using Dual-field arithmetic & Microcode-based architecture.

David Narh Amanor, Christof Paar, Jan Pelzl, "Efficient Hardware Architectures For Modular Muliplication On FPGAs" The first problem is the comparison of P and M. In the worst case all bits Of M and P must be compared. This problem can be solved by an approximative comparison with 2n instead of M.

The second problem is the number of additions or subtractions. All these operations can be replaced by one addition only. In order to do so, we pre estimate the number of Ms to be subtracted, and find out if Y will be added in the next iteration of the loop. There are only few possible values for this estimation.

III. IMPLEMENTING PRIME FIELD AND BINARY EXTENSION FIELD IN FPGA

The Prime field and Binary extension field consist of several hardware blocks. In this section we discuss about the hardware design involvement of both prime and Binary extension field.

A.PRIME FIELD UNIT ARCHITECTURE

Blakley introduced an algorithm to perform Modular multiplication of two integers A and B modulo an integer. It is an iterative binary double-and-add algorithm. The main idea of the algorithm is that it keeps the intermediate result after each iteration below the modulus value, which it avoids final division. In this paper, the modulus M corresponds to p and we say it Fp multiplication. All arithmetic in Fp are performed in two's complement number system, which avoids input and output conversions like existing implementations. The main difficulty of the Blakley algorithm is the computation of addition on large operands. The modified Blakley algorithm for large operands. The use of carry save adder (CSA) helps to speed up the repeated additions on large operands. However these modified versions require at least one final addition on large carry chain. Some pre-computed values too are used by this technique which requires additional time and storage area.

B.BINARY EXTENSION FIELD UNIT

Blakley introduced an algorithm to perform Modular multiplication of two integers A and B modulo an integer M. It is an iterative binary double-and-add algorithm. The main idea of the algorithm is that it keeps the intermediate result after each iteration below the modulus value, which it avoids final division. In this paper, the modulus M corresponds to p and we say it Fp multiplication. All arithmetic in Fp Blakley introduced an algorithm to perform Modular multiplication of two integers



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A and B modulo an integer M. It is an iterative binary double-and-add algorithm. The main idea of the algorithm is that it keeps the intermediate result after each iteration below the modulus value, which it avoids final division. In this paper, the modulus M corresponds to p and we say it Fp multiplication. All arithmetic in Fp are performed in two's complement number system, which avoids input and output conversions like existing implementations [6], [7]. The main difficulty of the Blakley algorithm is the computation of addition on large operands. The modified Blakley algorithm for large operands is shown in [8] and [9]. The use of carry save adder (CSA) helps to speed up the repeated additions on large operands. However these modified versions require at least one final addition on large carry chain. Some pre-computed values too are used by this technique which requires additional time and storage area.

IV. EXECUTION OF PARALLEL INSTRUCTIONS

The binary extension unit designed for the polynomial equation $\mathbf{p}(\mathbf{x})=\mathbf{x}\mathbf{4}+\mathbf{x}+\mathbf{1}$. Therefore out of 256 input bits 4 bits are taken for the each binary extension unit blocks. The binary extension field unit consists of GF addition, multiplication and double. The binary field addition is designed by using the XOR gate operation. The multiplier unit is designed by the look up table procedure. So this block is consist of memory blocks, which are used for storing the all possible output. The GF double unit consists of a simple addition unit.

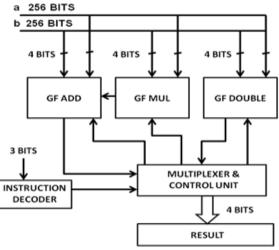


Fig 1. Quad core structure block diagram

The execution unit is named as configurable arithmetic unit (CAU). The CAU unit contains the both blocks of Prime field instruction execution unit and Binary extension field execution unit. The data accesses and instruction sequences are hard coded into the *sequence control* of the architecture which avoids the additional software development costs. The quad core contains the 4 blocks of CAU unit, each one will execute any one of prime field or binary extension field instructions and produce the result.



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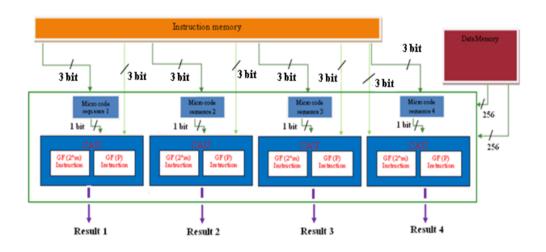


Fig 2. Block Diagram of Binary Extension field

1.000,999,995,99		0,999,997 ps UUUU 1110 00 01 10 11 00011011 1111 0000	1,000,993,996 ps	1,000,999,999 ps	1,001,000,000 p
1		0000 1110 00 01 10 11 00011011 1111			
		00 01 10 11 00011011 1111			
		01 10 11 00011011 1111			
		10 11 00011011 1111			
		11 00011011 1111			
		00011011			
		1111			X
					X
		0000			
		0000)
00000	0000	000000000000			X
00000	0000	000000000000)
		0			X
		0			
		0			X
		0			
		0			
10000	1110	000100000001			
10000	.1110	000100000001			
	.0000	.0000	0 0 0 0000 0 0000 0 0000 0 0 0 0 0 0 0	0 0 0000 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0

V. SIMULATION RESULT

Fig 3.Simuation Result

Various combination of instruction execution in parallel is verified. The instructions selection and input a and b value loading is in build in the coding. Example execution for Prime field addition, subtraction and binary extension field multiplication and double for the same input value of a=5h, b=4h, and p=7h is executed and the simulation output result is shown in the Fig. The output for the prime field addition is 102h, subtraction is 1h, Binary extension field multiplication is 7h, and binary extension field double is 1h.



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VI. RTL SCHEMATIC DIAGRAM

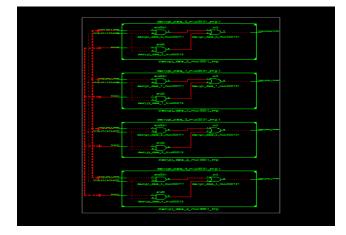


Fig 4. RTL Schematic Diagram

VII. HARDWARE MODEL

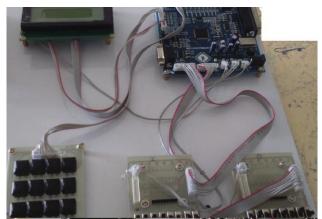


Fig 5.Hardware model

FPGA is a semiconductor device that based around a matrix of configurable logic blocks(CLB) connected through programmable interconnects. FPGA can be programmed to the desired application (or) functionality requirements although one-time programmable(OTP) FPGAs are available.

- A plane and regular structure in which logic and interconnect both are programmable.
- Programmability of logic -any combinational or sequential logic can be implemented.
- Programmability of interconnect –any logic component can be connected to anyone else.

Programmable interconnects

• Connection box: Connects input/output of logic block to interconnect channels



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- Switch box: Connects horizontal channels to vertical channels
- Transmission gate (or a pass transistor) is used for each condition.

VIII. CONCLUSION

The quad core structure executes the parallelism with both fields of instructions. At the same time 4 numbers of instructions are executing, this parallelism action will increase the speed. This architecture is an example design of crypto processor for executing prime field and binary field operations. This work has been further developed for the superscalar architecture with more number of instructions.

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